

QBF- and SAT-Based Synthesis from Safety Specs

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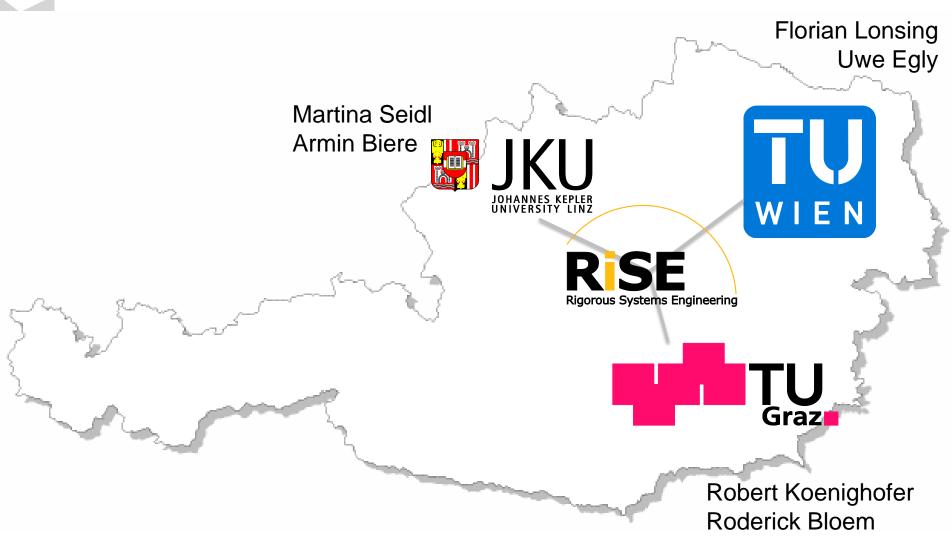
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IIAIK





Rich-Model Toolkit / RiSE Collaboration

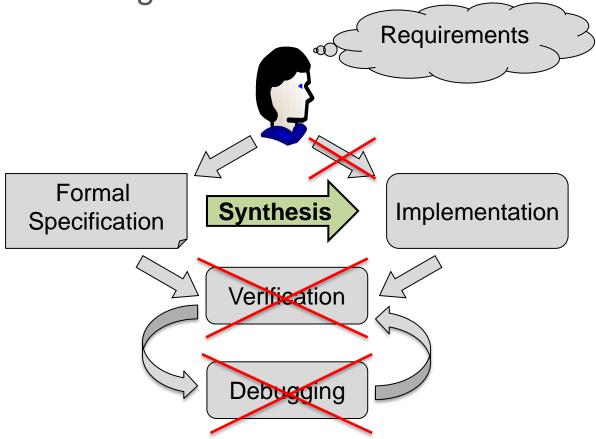






Motivation: Synthesis

Typical design flow:

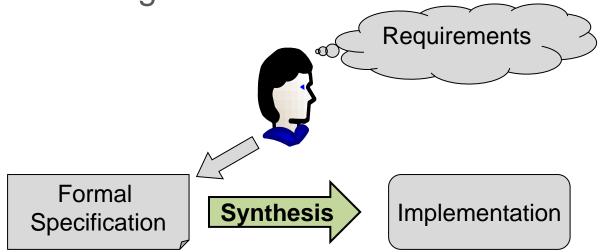




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Motivation: Synthesis

Typical design flow:







Motivation: From BDDs to SAT

- Challenge: scalability
 - → Symbolic algorithms
 - → Often implemented with BDDs
 - Known scalability issues
- Enormous achievements in decision procedures
 - SAT-solver, QBF-solvers, EPR-solvers, ...
 - → Exploit for synthesis





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Outline

- Problem definition
- Learning-based synthesis method
- Template-based synthesis method
- Extensions
- Experimental results



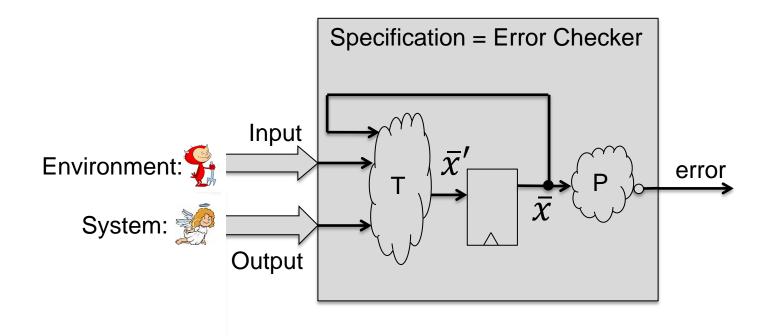


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Problem:

Synthesis from Safety Specifications

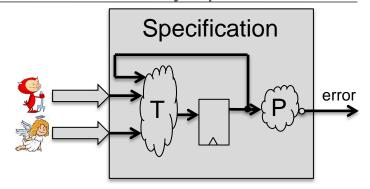
- "Something bad must never happen"
- Format:







Typical Synthesis Flow



- 1. Compute game graph
- 2. Compute "Winning Region" W
 - Set of states from which the system gran win
 - No matter what the environment does
 - Safety: ... stay in safe states
- 3. Compute a strategy
 - What to do in which situation in order to win
 - Safety: stay in winning region
- 4. Output strategy
 - E.g., as Verilog circuit



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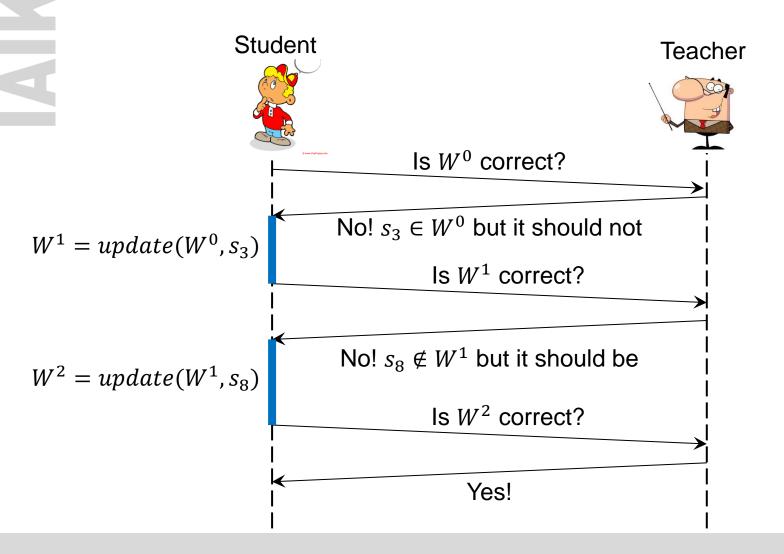
Learning-Based Synthesis Method







Supervised Learning







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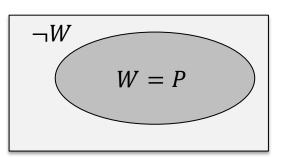
- $Force^{e}(A)$
 - the environment \(\frac{1}{2} \) can enforce to reach A in one step





- $Force^{e}(A)$:
 - the environment \(\frac{1}{2} \) can enforce to reach A in one step

$$W := P$$

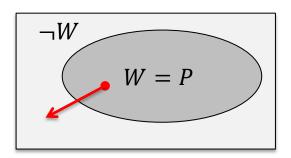






- $Force^{e}(A)$:
 - the environment \(\frac{1}{2} \) can enforce to reach A in one step

```
W \coloneqq P
while(sat(W \land Force^e(\neg W))) {
   pick s \vDash W \land Force^e(\neg W)
}
```



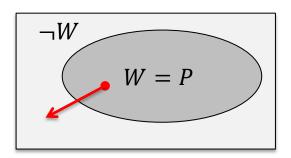




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```
W \coloneqq P
while(sat(W \land Force^e(\neg W))) {
   pick s \vDash W \land Force^e(\neg W)

W \coloneqq W \land \neg s
}
```



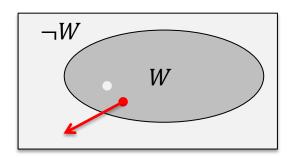




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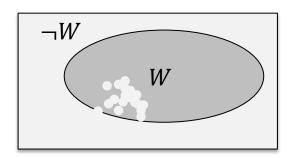




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W \coloneqq P
while(sat(W \land Force^e(\neg W))) {
   pick s \vDash W \land Force^e(\neg W)

W \coloneqq W \land \neg s
}
```

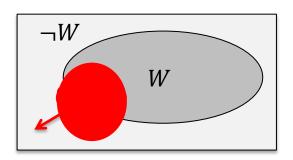




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- $Force^{e}(A)$:
 - the environment \(\frac{1}{2} \) can enforce to reach A in one step

```
W \coloneqq P
while(sat(W \land Force^e(\neg W))) {
    pick s \vDash W \land Force^e(\neg W)
    s \coloneqq generalize(s)
    W \coloneqq W \land \neg s
}
```

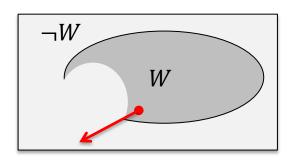






- $Force^{e}(A)$:
 - the environment \(\frac{1}{2} \) can enforce to reach A in one step

```
W \coloneqq P
while(sat(W \land Force^e(\neg W))) {
    pick s \vDash W \land Force^e(\neg W)
    s := generalize(s)
    W \coloneqq W \land \neg s
}
```

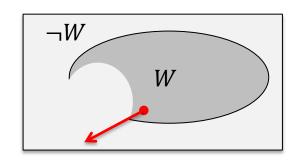






- $Force^{e}(A)$:
 - the environment \$\forall can enforce to reach A in one step

```
W \coloneqq P
\text{while}(\operatorname{sat}(W \land Force^{e}(\neg W))) \{
\text{pick } s \vDash W \land Force^{e}(\neg W)
\text{s} \coloneqq \text{generalize}(s)
W \coloneqq W \land \neg s
\}
\text{QBF}
\text{Solver}
\text{Satisfying}
x_{1} \land \neg x_{2} \land \neg x_{3} \land x_{4}
```







- $Force^{e}(A)$:
 - the environment \$\forall can enforce to reach A in one step

```
\begin{array}{c} W \coloneqq P \\ \text{while}(\operatorname{sat}(W \land Force^e(\neg W))) \{\\ \text{pick } s \vDash W \land Force^e(\neg W) \\ \text{s} \coloneqq \operatorname{generalize}(s) \\ W \coloneqq W \land \neg s \\ \} \end{array}
```





Template-Based Synthesis Method



Template-Based Method

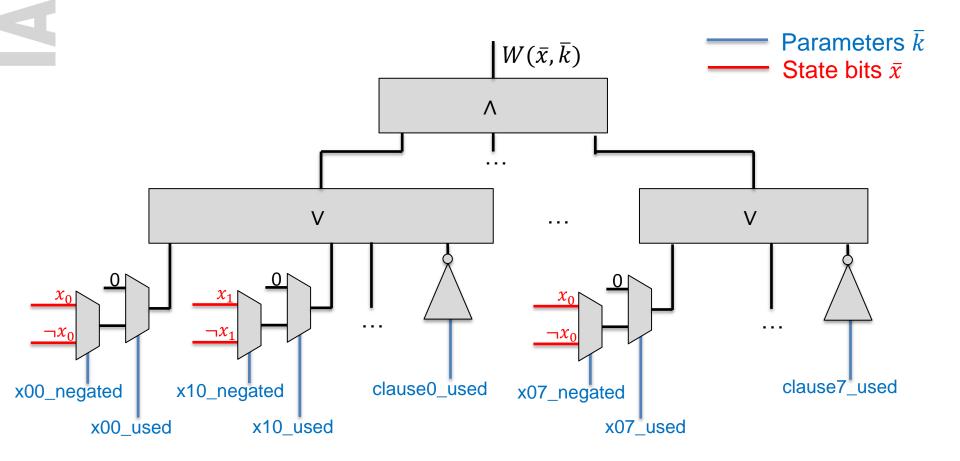
- Need to find $W(\bar{x})$ such that:
 - $I(\bar{x}) \to W(\bar{x})$
 - $W(\bar{x}) \to P(\bar{x})$
 - $W(\bar{x}) \to Force^{s}(W(\bar{x}))$
- Let $W(\bar{x}, \bar{k})$ be a parameterized function
 - Concrete values for $\bar{k} \rightarrow \text{concrete function } W(\bar{x})$
- Solve: $\exists \overline{k} \colon I(\overline{x}) \to W(\overline{x}, \overline{k}) \land W(\overline{x}, \overline{k}) \to P(\overline{x}) \land W(\overline{x}, \overline{k}) \to Force^s \left(W(\overline{x}, \overline{k})\right)$



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Template-Based Method:

CNF Template





Extensions

Templates and learning:

- QBF: Pre-processing
 - Extension of Blogger to preserve models

Learning-based method:

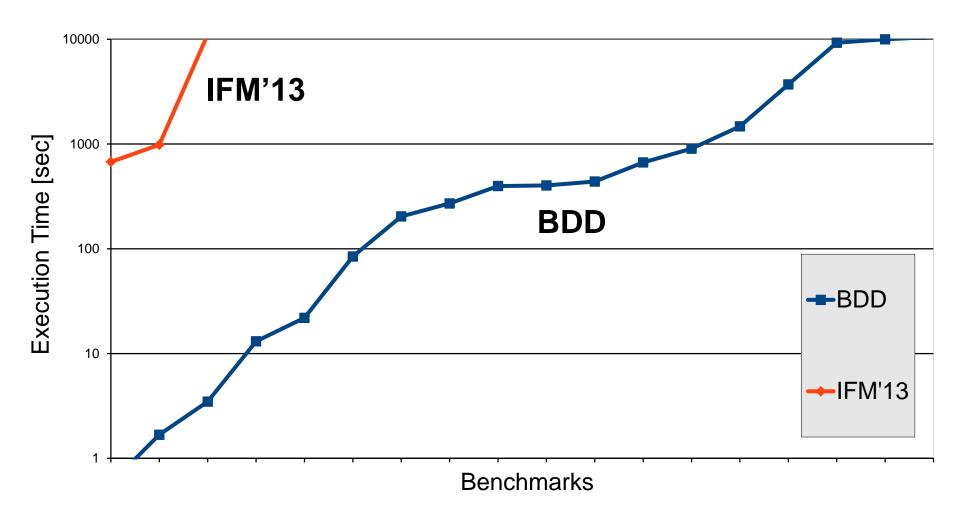
- SAT-based implementation
- Parallelized implementation



Experimental Results

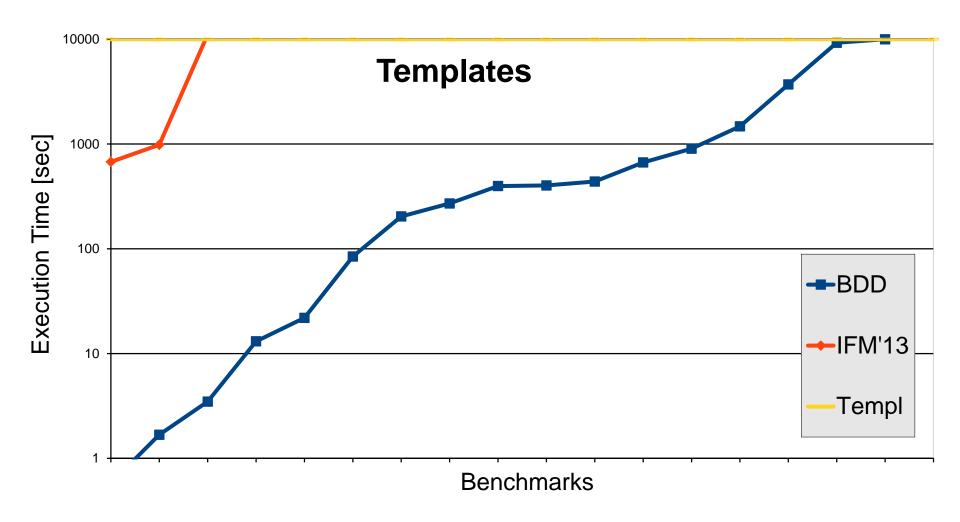






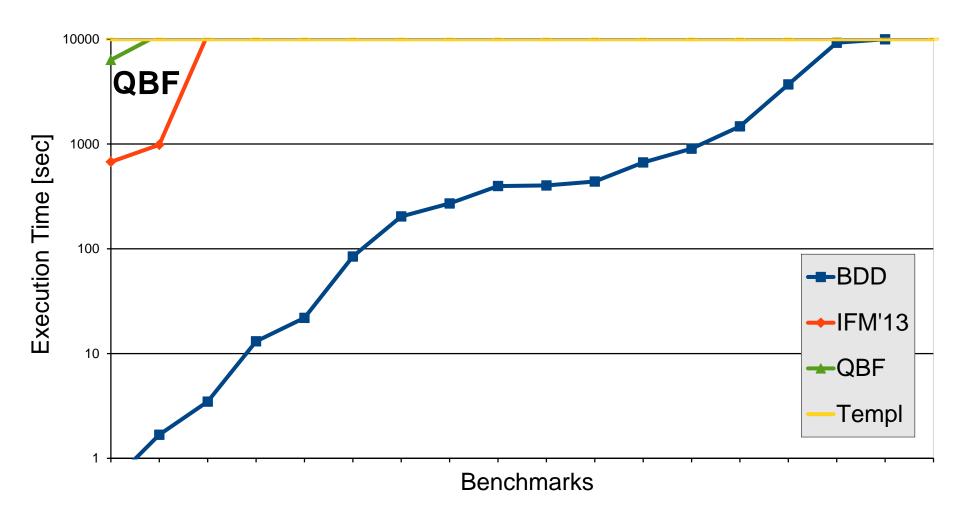






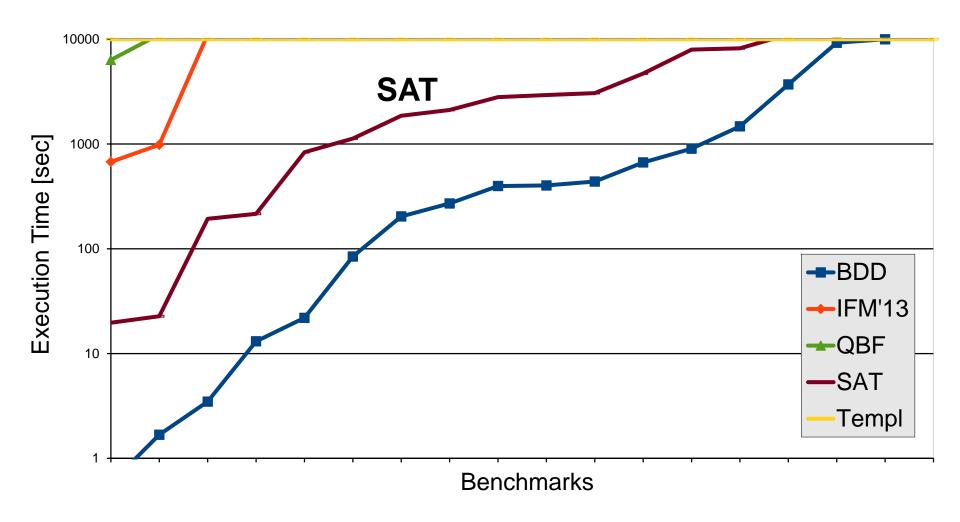






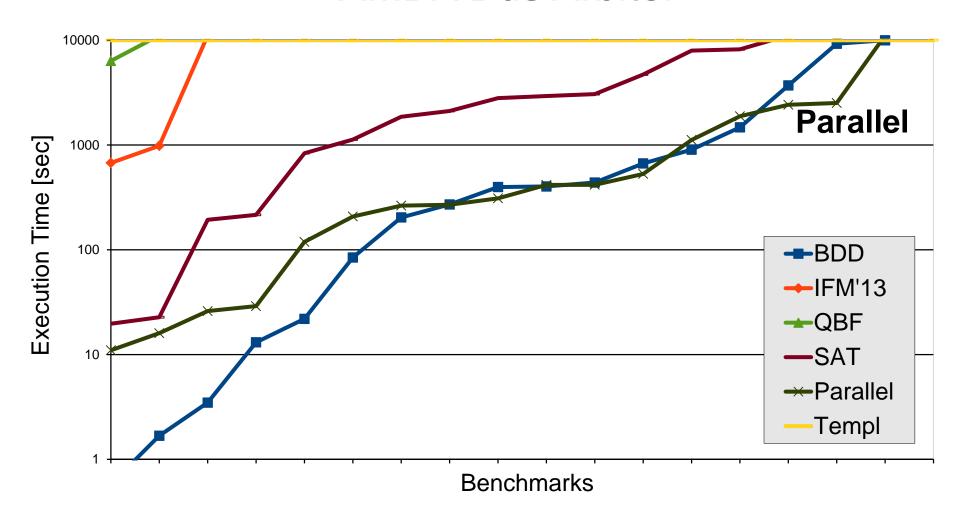








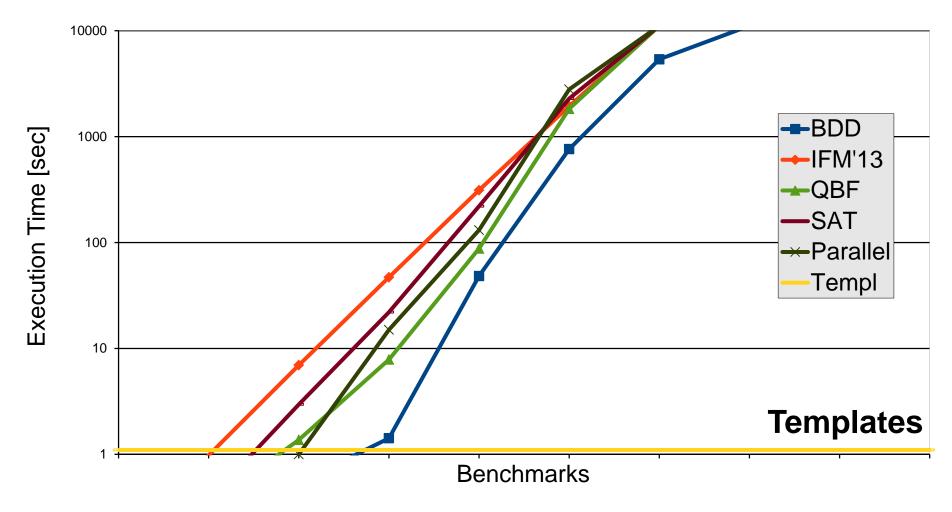








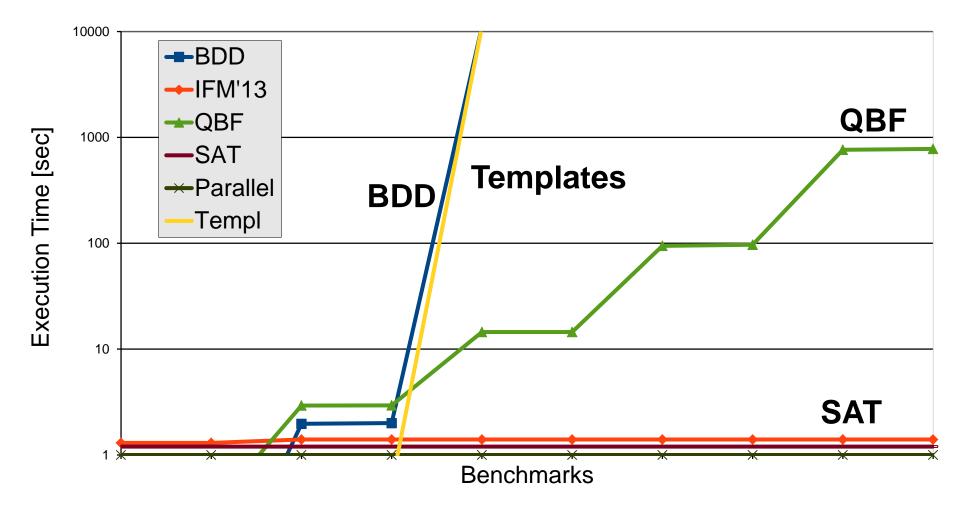
First Experiments: Combinational Multiplier





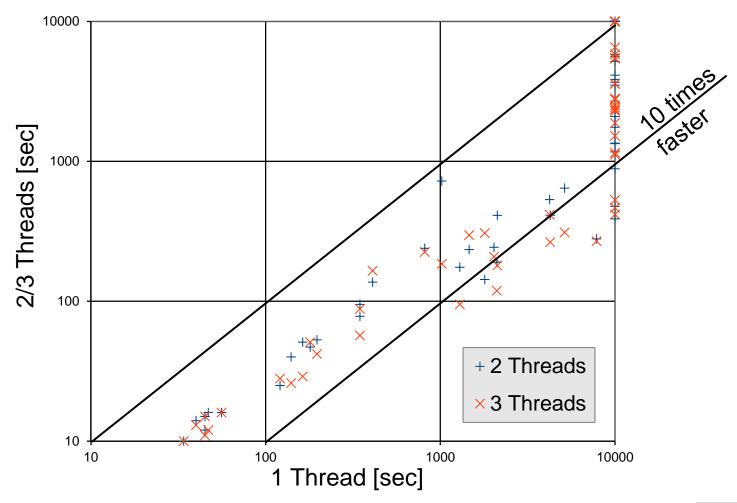


First Experiments: Barrel Shifter





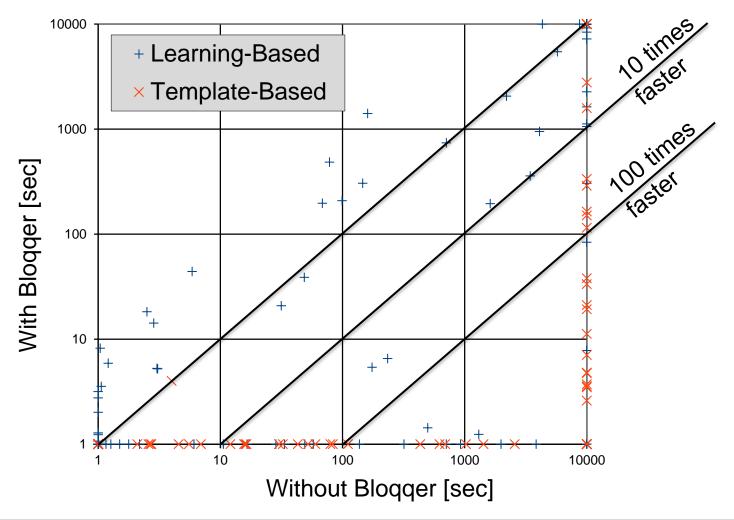
Parallelization Speedup







QBF Preprocessing Speedup:







Conclusions

- No clear winner
 - Different methods are good at different benchmarks
- SAT-based implementation faster than QBF
 - Room for optimization in QBF
- Parallelization is beneficial
 - Different solvers complement each other
- Tool:
 - Open-source release in progress
 - http://www.iaik.tugraz.at/content/research/design_verification/demiurge/

