# Automata on Infinite Trees

# Büchi Automata on Infinite Trees

## Definition

A Büchi tree automaton over $\Sigma$ is $A = \langle S, I, T, F \rangle$, where:

- $S$ is a finite set of *states*,

- $I \subseteq S$ is a set of *initial states*,

- $T : S \times \Sigma \to 2^{S \times S}$ is the *transition relation*,

- $F \subseteq S$ is the set of *final states*.

# Runs

A *run* of $A$ over a tree $t : \{0,1\}^* \to \Sigma$ is a mapping $\pi : \{0,1\}^* \to S$ such that, for each position $p \in \{0,1\}^*$, where $q = \pi(p)$, we have:

- if $p = \epsilon$ then $q \in I$, and

- if $q_i = \pi(pi)$, $i = 0, 1$ then $\langle q_0, q_1 \rangle \in T(q, t(p))$.

If $\pi$ is a run of $A$ and $\sigma$ is a path in $t$, let $\pi_{|\sigma}$ denote the path in $\pi$ corresponding to $\sigma$.

A run $\pi$ is said to be *accepting*, if and only if for every path $\sigma$ in $t$ we have:

$$\inf(\pi_{|\sigma}) \cap F \neq \emptyset$$

# Closure Properties

For every Büchi automaton $A$ there exists a complete Büchi automaton $A'$ such that $\mathcal{L}(A) = \mathcal{L}(A')$.

**Theorem 1** *The class of Büchi-recognizable tree languages is closed under union, intersection and projection.*

Let $A_i = \langle S_i, I_i, T_i, F_i \rangle$, $i = 1, 2$, where $S_1 \cap S_2 = \emptyset$.

Let $A_\cup = \langle S_1 \cup S_2, I_1 \cup I_2, T_1 \cup T_2, F_1 \cup F_2 \rangle$.

## Closure Properties

Let $A_\cap = \langle S, I, T, F \rangle$ where:

- $S = S_1 \times S_2 \times \{0, 1, 2\}$

- $I = I_1 \times I_2 \times \{1\}$

- for any $s, s_1, s_2 \in S_1$, $s', s_1', s_2' \in S_2$, $a, b \in \{0, 1, 2\}$:

$$\langle (s_1, s_1', b), (s_2, s_2', b) \rangle \in T((s, s', a), \sigma)$$

  iff $\langle s_1, s_2 \rangle \in T(s, \sigma)$, $\langle s_1', s_2' \rangle \in T(s', \sigma)$ and:

  1. if $a = 0$ or ($a = 1$ and $s \notin F_1$), then $b = 1$
  2. if ($a = 1$ and $s \in F_1$) or ($a = 2$ and $s \notin F_1$), then $b = 2$
  3. if $a = 2$ and $s' \in F_2$, then $b = 0$

- $F = S \times S \times \{0\}$

# Emptiness of Büchi Tree Automata

Let $A = \langle S, I, T, F \rangle$ be a Büchi tree automaton where $F = \{s_1, \ldots, s_m\}$, and $\pi : \{0, 1\}^* \to S$ be a successful run of $A$ on the tree $t \in \mathcal{T}(\Sigma)$.

For any $s \in S$, and any $u \in \{0, 1\}^*$ such that $\pi(u) = s$, let

$$d_u^\pi = \{w \in u \cdot \{0, 1\}^* \mid \pi(v) \notin F, \text{ for all } u < v < w\}$$

By König's lemma, $d_u^\pi$ is finite for any $u \in \{0, 1\}^*$.

Let $t_s^\pi$ be the restriction of $t$ to $d_u^\pi$. Let

$$T_s = \{t_s^\pi \mid \pi \text{ is a successful run of } A \text{ on } t\}$$

# Emptiness of Büchi Tree Automata

If $\vec{s} = \langle s_1, \ldots, s_m \rangle$:

$$\mathcal{L}(A) = \bigcup_{s_0 \in I} T_{s_0} \cdot_{\vec{s}} \langle T_{s_1}, \ldots, T_{s_m} \rangle^{\omega \vec{s}}$$

Conversely, the expression above denotes a Büchi-recognizable tree language.

Let $A = \langle S, I, T, F \rangle$ be a Büchi tree automaton. For each $s \in S$ let $T_s$ be the recognizable tree language defined above. Eliminate from $S$ (and $T$) all states $s$ such that $T_s = \emptyset$, and let $S'$ be the resulting set of states.

We claim that $\mathcal{L}(A) \neq \emptyset \iff S' \cap I \neq \emptyset$.

# The Complement Problem

Let $\Sigma = \{a, b\}$, $\mathcal{T}_0 = \{t \in \mathcal{T}^\omega(\Sigma) \mid$ some path in $t$ has infinitely many $a$'s$\}$

$\mathcal{T}_0$ is Büchi recognizable.

Let $A = \langle \{s_0, s_1, s_a, s_b\}, \{s_0\}, T, \{s_1, s_a\} \rangle$, where $T$ is defined by:

$$
\begin{aligned}
a(s_{0,a,b}) &\rightarrow \{\langle s_1, s_a \rangle, \langle s_a, s_1 \rangle\} \\
b(s_{0,a,b}) &\rightarrow \{\langle s_1, s_b \rangle, \langle s_b, s_1 \rangle\} \\
a(s_1) &\rightarrow \{\langle s_1, s_1 \rangle\} \\
b(s_1) &\rightarrow \{\langle s_1, s_1 \rangle\}
\end{aligned}
$$

# The Complement Problem

Let $\mathcal{T}_1 = \mathcal{T}^\omega(\Sigma) \setminus \mathcal{T}_0 = \{t \in \mathcal{T}^\omega(\Sigma) \mid$ all paths in $t$ have finitely many $a$'s$\}$. We show that $\mathcal{T}_1$ cannot be recognized by a Büchi tree automaton.

***Exercise 1*** $I = \{s_0, s_1\}$, $F = \{s_1\}$ *and*

$$
\begin{aligned}
a(s_0) &\rightarrow \langle s_0, s_0 \rangle \\
&\phantom{\rightarrow} \langle s_0, s_1 \rangle \\
&\phantom{\rightarrow} \langle s_1, s_0 \rangle \\
&\phantom{\rightarrow} \langle s_1, s_1 \rangle \\
b(s_0) &\rightarrow \langle s_0, s_0 \rangle \\
&\phantom{\rightarrow} \langle s_0, s_1 \rangle \\
&\phantom{\rightarrow} \langle s_1, s_0 \rangle \\
&\phantom{\rightarrow} \langle s_1, s_1 \rangle \\
b(s_1) &\rightarrow \langle s_1, s_1 \rangle
\end{aligned}
$$

# The Complement Problem

Let $T_n : \{0,1\}^* \to \Sigma$ be the language of trees:

$$t_n(p) = \begin{cases} \text{a} & \text{if } p \in \{\epsilon, 1^{m_1}0, 1^{m_1}01^{m_2}0, \ldots, 1^{m_1}01^{m_2}0 \ldots 1^{m_n}0 \mid m_1, \ldots m_n \in \mathbb{N}\} \\ \text{b} & \text{otherwise} \end{cases}$$

Obviously, $T_n \subset \mathcal{T}_1$, for all $n \in \mathbb{N}$.

Suppose there exists a Büchi automaton $A = \langle S, I, T, F \rangle$ with $k$ states, s.t. $\mathcal{L}(A) = \mathcal{T}_1$. Let $\pi$ be the accepting run of $A$ over $t_{k+1}$. Then there exist:

- $m_1 > 0$ such that $\pi(1^{m_1}) = s_1 \in F$

- $m_2 > 0$ such that $\pi(1^{m_1}01^{m_2}) = s_2 \in F$

- $\ldots$

There exists a path $\sigma$ in $t_m$ and $u < v < w < \sigma$, such that $\pi(u) = \pi(w) = s \in F$ and $t_m(v) = a$. Then $\pi = r_1 \cdot_s r_2 \cdot_s r_3$, and $r_1 \cdot_s r_2^{\omega s}$ is a successful run on $q_1 \cdot q_2^\omega$, which contains a path with infinitely many $a$.

# Müller Automata on Infinite Trees

# Definition

A Müller tree automaton $\Sigma$ is $A = \langle S, I, T, \mathcal{F} \rangle$, where:

- $S$ is a finite set of *states*,

- $I \subseteq S$ is a set of *initial states*,

- $T : S \times \Sigma \to 2^{S \times S}$ is the *transition function*,

- $\mathcal{F} \subseteq 2^S$, is the set of *accepting sets*.

A run $\pi$ of $A$ over $t$ is said to be *accepting*, iff for every path $\sigma$ in $t$:

$$\mathrm{inf}(\pi_{|\sigma}) \in \mathcal{F}$$

# Closure Properties

The class of Müller-recognizable tree languages is closed under union and intersection.

For union, the proof is exactly as in the case of Büchi automata. For $A_\cup$, the set of accepting sets is the union of the sets $\mathcal{F}_i$, $i = 1, 2$.

For intersection, let $A_\cap = \langle S_1 \times S_2, I_1 \times I_2, T, \mathcal{F} \rangle$, where:

- $\langle (s_1, s_1'), (s_2, s_2') \rangle \in T((s, s'), \sigma)$ iff $\langle s_1, s_2 \rangle \in T(s, \sigma)$ and $\langle s_1', s_2' \rangle \in T(s', \sigma)$, and

- $\mathcal{F} = \{G \in S_1 \times S_2 \mid pr_1(G) \in \mathcal{F}_1 \text{ and } pr_2(G) \in \mathcal{F}_2\}$, where:
    - $pr_1(G) = \{s \in S_1 \mid \exists s' \, . \, (s, s') \in G\}$, and
    - $pr_2(G) = \{s \in S_2 \mid \exists s' \, . \, (s', s) \in G\}$.

# Rabin Automata on Infinite Trees

# Definition

A Rabin tree automaton $\Sigma$ is $A = \langle S, I, T, \Omega \rangle$, where:

- $S$ is a finite set of *states*,

- $I \subseteq S$ is a set of *initial states*,

- $T : S \times \Sigma \to 2^{S \times S}$ is the *transition function*,

- $\Omega = \{\langle N_1, P_1 \rangle, \ldots, \langle P_n, N_n \rangle\}$ is the set of *accepting pairs*.

A run $\pi$ of $A$ over $t$ is said to be *accepting*, if and only if for every path $\sigma$ in $t$ there exists a pair $\langle N_i, P_i \rangle \in \Omega$ such that:

$$\inf(\pi_{|\sigma}) \cap N_i = \emptyset \text{ and } \inf(\pi_{|\sigma}) \cap P_i \neq \emptyset$$

# Büchi, Müller and Rabin

For every Büchi tree automaton $A$ there exists a Rabin tree automaton $B$, such that $\mathcal{L}(A) = \mathcal{L}(B)$, but not viceversa.

For every Müller tree automaton $A$ there exists a Rabin tree automaton $B$, such that $\mathcal{L}(A) = \mathcal{L}(B)$, and viceversa.

# The Rabin Complementation Theorem

**Theorem 2 (Rabin '69)** *The class of Rabin-recognizable tree languages is closed under complement.*

The class of Rabin-recognizable tree languages is closed under union and intersection.

# Emptiness of Rabin Automata

Given an alphabet $\Sigma$, an infinite tree $t \in \mathcal{T}^\omega(\Sigma)$ is said to be *regular* if there are only finitely many distinct subtrees $t_u$ of $t$, where $u \in \{0, 1\}^*$.

***Example 1*** *The infinite binary tree $f(g(f(\ldots), f(\ldots)), g(f(\ldots), f(\ldots)))$ is regular.* $\square$

## Theorem 3 (Rabin '72)

1. *Any non-empty Rabin-recognizable set of trees contains a regular tree.*

2. *The emptiness problem for Rabin tree automata is decidable.*

# Reduction to empty alphabet

Let $A = \langle S, I, T, \Omega \rangle$ be a Rabin tree automaton over $\Sigma$, such that $\mathcal{L}(A) \neq \emptyset$, where $\Omega = \{\langle N_1, P_1 \rangle, \ldots, \langle N_n, P_n \rangle\}$.

Let $A' = \langle S \times \Sigma, I \times \Sigma, T', \Omega' \rangle$, where:

- $\langle (s_1, \sigma_1), (s_2, \sigma_2) \rangle \in T'((s, \sigma))$ iff $\langle s_1, s_2 \rangle \in T(s, \sigma)$, and $\sigma_1, \sigma_2 \in \Sigma$.

- $\Omega' = \{\langle N_1 \times \Sigma, P_1 \times \Sigma \rangle, \ldots, \langle N_n \times \Sigma, P_n \times \Sigma \rangle\}$.

The successful runs of $A'$ are pairs $(\pi, t)$, where $t \in \mathcal{L}(A)$, and $\pi$ is a successful run of $A$ on $t$.

# Regular successful runs

For any Rabin tree automaton $A$, there exists a Rabin tree automaton $A'$ with one initial state such that $\mathcal{L}(A) = \mathcal{L}(A')$.

Consider a Rabin tree automaton $A = \langle S, s_0, T, \Omega \rangle$ over the empty alphabet, and let $\pi$ be a successful run of $A$.

**Claim 1** *If $A$ has a successful run, $A$ has also a regular successful run.*

A state $s \in S$ is said to be *live* if $s \neq s_0$ and $\langle s_1, s_2 \rangle \in T(s)$ for some $s_1, s_2 \in S$, where either $s_1 \neq s$ or $s_2 \neq s$.

By induction on $n =$ the number of live states in $A$.

## Regular successful runs

If $n = 0$, $\pi(\epsilon) = s_0$ and $\pi(p) = s$, for all $p \in dom(\pi)$, and $s \in S$ non-live.

**Case 1** If some live state in $A$ is missing on $\pi$, apply the induction hypothesis.

**Case 2** All states of $A$ appear on $\pi$, and there is a position $u \in \{0,1\}^*$ such that $\pi(u) = s$ is live, but some live state $s'$ does not appear in $\pi_u$.

Let $\pi_1 = \pi \setminus \pi_u$ and $\pi_2 = \pi_u$. Both $\pi_1$ and $\pi_2$ are runs of automata with $n-1$ live states, hence there exists successful regular runs $\pi_1'$ and $\pi_2'$ of these automata. The desired run is $\pi_1' \cdot_s \pi_2'$.

# Regular successful runs

**Case 3** All live states appear in any subtree of $\pi$. Let $\sigma$ be a path in $\pi$ consisting of all the live states appearing again and again, and only of the live states, with the exception of $\pi(\epsilon)$. **Q:** Why does $\sigma$ exist?

There exists $\langle N, P \rangle \in \Omega$, such that $\inf(\sigma) \cap N = \emptyset$ and $\inf(\sigma) \cap P \neq \emptyset$. Then $N$ contains only non-live states.

Let $s \in \inf(\sigma) \cap P$ and $u$, $v$ be the $1^{st}$ and $2^{nd}$ positions such that $\sigma(u) = \sigma(v) = s$.

Let $\pi_1 = \pi \setminus \pi_u$ and $\pi_2 = \pi_u \setminus \pi_v$. Both $\pi_1$ and $\pi_2$ are runs of automata with $n - 1$ live states, hence there exists successful regular runs $\pi_1'$ and $\pi_2'$ of these automata. The desired run is $\pi_1' \cdot_s {\pi_2'}^{\omega s}$.

# The Emptiness Problem

Let $A$ be an input-free Rabin tree automaton with $n$ live states.

We derive $A_{n-1}, A_{n-2}, \ldots, A_0$ from $A$, having $n-1, n-2, \ldots 0$ live states.

If $A$ has a successful run, then it it has a regular run, composed of runs of $A_{n-1}, A_{n-2}, \ldots, A_0$.

So it is enough to check emptiness of $A_{n-1}, A_{n-2}, \ldots, A_0$.

# Rabin Automata, SkS and S$\omega$S

# Defining infinite paths

We say that a set of positions $X$ is linear iff the following holds:

$$linear(X) \; : \; (\forall x, y \; . \; X(x) \wedge X(y) \rightarrow x \leq y \vee y \leq x)$$

$X$ is a path iff:

$$path(X) \; : \; linear(X) \wedge \forall Y \; . \; linear(Y) \wedge X \subseteq Y \rightarrow X = Y$$

# From Automata to Formulae

Let $A = \langle S, I, T, \Omega \rangle$ be a Rabin tree automaton, where $S = \{s_1, \ldots, s_p\}$.

Let $\vec{Y} = \{Y_1, \ldots, Y_p\}$ be set variables.

If $X$ denotes a path, state $i$ appears infinitely often in $X$ iff:

$$inf_i(X) \; : \; \forall x \, . \, X(x) \rightarrow \exists y \, . \, x \leq y \wedge X(y) \wedge Y_i(y)$$

The formula expressing the accepting condition is:

$$\Phi_\Omega(\vec{Y}) \; : \; \forall X \, . \, path(X) \rightarrow \bigvee_{\langle N, P \rangle \in \Omega} \left( \bigwedge_{s_i \in N} \neg inf_i(X) \; \wedge \; \bigvee_{s_i \in P} inf_i(X) \right)$$

# Decidability of S2S

**Theorem 4** *Given an alphabet $\Sigma$, a tree language $L \subseteq \mathcal{T}^\omega(\Sigma)$ is definable in S2S iff it is recognizable.*

**Corollary 1** *The SAT problem for S2S is decidable.*

# Obtaining Decidability Results by Reduction

Suppose we have a logic $\mathcal{L}$ interpreted over the domain $\mathcal{D}$, such that the following problem is decidable:

$$\text{for each formula } \varphi \text{ of } \mathcal{L} \text{ there exists } \mathfrak{m} \in \mathcal{D} \text{ such that } \mathfrak{m} \models \varphi$$

Then we can prove the same thing for another logic $\mathcal{L}'$ interpreted over $\mathcal{D}'$ iff there exists functions $\Delta : \mathcal{D}' \to \mathcal{D}$ and $\Lambda : \mathcal{L}' \to \mathcal{L}$ such that for all $\mathfrak{m}' \in D'$ and $\varphi' \in \mathcal{L}$ we have:

$$\mathfrak{m}' \models \varphi' \iff \Delta(\mathfrak{m}') \models \Lambda(\varphi')$$

# Decidability of S$\omega$S

Every tree $t : \mathbb{N}^* \to \Sigma$ can be encoded as $t' : \{0,1\}^* \to \Sigma$.

Let $D = \{\epsilon\} \cup \bigcup_{n_1,\ldots,n_k \in \mathbb{N}} \{01^{n_1}01^{n_2}0\ldots 1^{n_k}0 \mid k \geq 1, 1 \leq i \leq k\}$.

Embedding the domain of S$\omega$S into S2S:

$$D(x) \quad : \quad \exists z \forall y \, . \, (z \leq y) \wedge x = z \ \vee$$

$$s_0(z) \leq x \wedge \forall y \, . \, z < y \wedge s_0(y) \leq x \to \exists y' \, . \, y = s_1(y')$$

# Decidability of S$\omega$S

If $p = 01^{n_1}01^{n_2}0\ldots1^{n_k}0$, let $f_i(p) = 01^{n_1}01^{n_2}0\ldots1^{n_k}01^i0$

$$x \preceq_1 y \quad : \quad D(x) \wedge D(y) \wedge x \preceq y$$

Define the relation $x \leq_D y$ iff $x \in D$ and $y = x01^n$, for some $n \in \mathbb{N}$.

Define $f_0, f_1, f_2, \ldots$ by induction:

- $f_0(x) = y \quad : \quad D(x) \wedge D(y) \wedge x \leq_D y \wedge \forall z \,.\, x \leq_D z \rightarrow y \preceq_1 z$

- $f_{i+1}(x) = y \quad : \quad D(x) \wedge D(y) \wedge x \leq_D y \wedge \forall z \,.\, x \leq_D z \wedge \bigwedge_{0 \leq k \leq i} z \neq f_k(x) \rightarrow y \preceq_1 z$