# Analysis and Verification of Higher Order Functional Programs: Automata-Theoretic Approach 

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# Automata-theoretic approach 

- Only requires reachability and fair termination verifiers
- A standard method for imperative programs


# Automata-theoretic approach applied to temporal verification of Ocaml programs with high-order procedures 

## Our work

- Monitoring for evaluation trees
- Product construction
- Evaluation


## Our work

- Monitoring for evaluation trees


## Product construction

## Evaluation

## A simple program

$1+2$

## Evaluation judgement

$\emptyset \vdash 1+2 \Rightarrow 3$

## Evaluation tree

$$
\begin{array}{ccc}
\hline \emptyset \vdash 2 \Rightarrow 2 & \emptyset \vdash 1 \Rightarrow 1 & \emptyset \vdash+\Rightarrow+ \\
& \emptyset \vdash 1+2 \Rightarrow 3
\end{array}
$$

## Evaluation of $1+2$

$\emptyset \vdash 1+2 \Rightarrow$

## Evaluation of 1+2

## $\emptyset \vdash 2 \Rightarrow$

$$
\emptyset \vdash 1+2 \Rightarrow
$$

## Evaluation of 1+2

$\emptyset \vdash 2 \Rightarrow 2$
$\emptyset \vdash 1+2 \Rightarrow$

## Evaluation of $1+2$

$$
\begin{array}{ll}
\hline \emptyset \vdash 2 \Rightarrow 2 & \emptyset \vdash 1 \Rightarrow \\
& \emptyset \vdash 1+2 \Rightarrow
\end{array}
$$

## Evaluation of 1+2

$$
\begin{aligned}
& \hline \emptyset \vdash 2 \Rightarrow 2 \emptyset \vdash 1 \Rightarrow 1 \\
& \emptyset \vdash 1+2 \Rightarrow
\end{aligned}
$$

## Evaluation of 1+2

$$
\begin{gathered}
\emptyset \vdash 2 \Rightarrow 2 \\
\emptyset \vdash 1 \Rightarrow 1 \\
\emptyset \vdash 1+2 \Rightarrow
\end{gathered}
$$

## Evaluation of 1+2

$$
\begin{array}{lll} 
& \emptyset \vdash 1 \Rightarrow 1 & \emptyset \vdash+\Rightarrow+ \\
& \emptyset \vdash 1+2 \Rightarrow
\end{array}
$$

## Evaluation of 1+2

$$
\begin{array}{ccc}
\hline \emptyset \vdash 2 \Rightarrow 2 & \emptyset \vdash 1 \Rightarrow 1 & \emptyset \vdash+\Rightarrow+ \\
& \emptyset \vdash 1+2 \Rightarrow 3
\end{array}
$$

# A simple monitor 

## Count additions

$1+2$

## Monitored tree

$$
\begin{array}{ccc}
\overline{\sigma_{2}^{\uparrow} \uparrow \emptyset \vdash 2 \Rightarrow 2 \downarrow \sigma_{2}^{\downarrow}} & \overline{\sigma_{3}^{\uparrow} \uparrow \emptyset \vdash 1 \Rightarrow 1 \downarrow \sigma_{3}^{\downarrow}} & \\
\sigma_{1}^{\uparrow \uparrow \emptyset \vdash 1+2 \Rightarrow 3 \downarrow \sigma_{1}^{\downarrow} \uparrow+\Rightarrow+\downarrow \sigma_{4}^{\downarrow}} \\
\hline
\end{array}
$$

# Monitored tree 

## Initial state: 0

$$
\begin{gathered}
\overline{\sigma_{2}^{\uparrow} \uparrow \emptyset \vdash 2 \Rightarrow 2 \downarrow \sigma_{2}^{\downarrow}} \quad \overline{\sigma_{3}^{\uparrow} \uparrow \emptyset \vdash 1 \Rightarrow 1 \downarrow \sigma_{3}^{\downarrow}} \quad \\
\sigma_{1}^{\uparrow \uparrow \emptyset \vdash 1+2 \Rightarrow 3 \downarrow \sigma_{1}^{\downarrow}}
\end{gathered}
$$

## Monitoring of I+2

## Initial state: $0<$ Current state

$$
\sigma_{2}^{\uparrow \uparrow \emptyset \vdash 2 \Rightarrow 2 \downarrow \sigma_{2}^{\downarrow}} \quad \overline{\sigma_{3}^{\uparrow} \uparrow \emptyset \vdash 1 \Rightarrow 1 \downarrow \sigma_{3}^{\downarrow}} \quad \sigma_{4}^{\uparrow \uparrow \emptyset \vdash+\Rightarrow+\downarrow \sigma_{4}^{\downarrow}}
$$

## Monitoring of I+2

## Initial state: $0<$ Current state

$\sigma_{2}^{\uparrow \uparrow \emptyset \vdash 2 \Rightarrow 2 \downarrow \sigma_{2}^{\downarrow}}$
$\overline{\sigma_{3}^{\uparrow} \uparrow \emptyset \vdash 1 \Rightarrow 1 \downarrow \sigma_{3}^{\downarrow}}$
$\sigma_{4}^{\uparrow \uparrow \emptyset \vdash+\Rightarrow+\downarrow \sigma_{4}^{\downarrow}}$
(0) $\uparrow \emptyset \vdash 1+2 \Rightarrow 3 \downarrow \sigma_{1}^{\downarrow}$

## Monitoring of $1+2$

$$
\begin{aligned}
\overline{\sigma_{2}^{\uparrow \uparrow \emptyset \vdash 2 \Rightarrow 2 \downarrow \sigma_{2}^{\downarrow}}} \quad \frac{\sigma_{3}^{\uparrow} \uparrow \emptyset \vdash 1 \Rightarrow 1 \downarrow \sigma_{3}^{\downarrow}}{} \quad & \begin{array}{l}
\sigma_{4}^{\uparrow \uparrow \emptyset \vdash+\Rightarrow+\downarrow \sigma_{4}^{\downarrow}} \\
\\
\\
\\
\text { Current state }
\end{array}
\end{aligned}
$$

## Monitoring of I+2



## Monitoring of I+2



## Monitoring of $1+2$

| $0 \uparrow \emptyset \vdash 2 \Rightarrow 2 \downarrow 0$ | (0) $\uparrow \emptyset \vdash 1 \Rightarrow 1 \downarrow \sigma_{3}^{\downarrow}$ | $\sigma_{4}^{\uparrow} \uparrow \emptyset \vdash+\Rightarrow+\downarrow \sigma_{4}^{\downarrow}$ |
| :---: | :---: | :---: |
|  | $0 \uparrow \emptyset \vdash 1+2 \Rightarrow 3 \downarrow$ ¢ |  |

## Monitoring of $1+2$

| $\overline{0 \uparrow \emptyset \vdash 2 \Rightarrow 2 \downarrow 0}$ | $\overline{0 \uparrow \emptyset \vdash 1 \Rightarrow 1 \downarrow 0}$ | $\overline{\sigma_{4}^{\uparrow} \uparrow \emptyset \vdash+\Rightarrow+\downarrow \sigma_{4}^{\downarrow}}$ |
| :---: | :---: | :---: |
| $0 \uparrow \emptyset \vdash 1+2 \Rightarrow 3 \downarrow \sigma_{1}^{\downarrow}$ |  |  |

## Monitoring of I+2

| $0 \uparrow \emptyset \vdash 2 \Rightarrow 2 \downarrow 0$ | $0 \uparrow \emptyset \vdash 1 \Rightarrow 1 \downarrow 0$ | (0) $\uparrow$ Ø $++\Rightarrow+\downarrow \sigma_{4}^{\downarrow}$ |
| :---: | :---: | :---: |
|  | $\uparrow \emptyset \vdash 1+2 \Rightarrow 3 \downarrow \sigma$ |  |

## Monitoring of I+2

## Current state

## Monitoring of I+2

## Current state

| $\overline{0 \uparrow \emptyset \vdash 2 \Rightarrow 2 \downarrow 0}$ | $\overline{0 \uparrow \emptyset \vdash 1 \Rightarrow 1 \downarrow 0} \quad$ | $\overline{0 \uparrow \emptyset \vdash+\Rightarrow+\downarrow 0}$ |
| :---: | :---: | :---: |
| $0 \uparrow \emptyset \vdash 1+2 \Rightarrow 3 \downarrow(1)$ |  |  |

# Addition counting monitor 

$$
\mathrm{M}_{+}=\left(\mathbb{N}, 0, \rho_{+}\right)
$$

# Addition counting monitor 

$$
\mathrm{M}_{+}=\left(\mathbb{N}, 0, \rho_{+}\right)
$$

## Monitor states

Addition counting monitor

$$
\mathrm{M}_{+}=\left(\mathbb{N}, 0, \rho_{+}\right)
$$

Monitor states

Initial state

Addition counting

## monitor

$$
\mathrm{M}_{+}=\left(\mathbb{N}, 0, \rho_{+}\right)
$$



## State transition function

$$
\mathrm{M}_{+}=\left(\mathbb{N}, 0, \rho_{+}\right)
$$

$$
\rho_{+}(\sigma, j, \delta)=\left\{\begin{array}{lc}
\sigma+1 & \text { if } \delta=\downarrow \wedge \\
& \exists \mathcal{E}, e_{1}, e_{2}, v: j=\mathcal{E} \vdash e_{1}+e_{2} \Rightarrow v \\
\sigma & \text { otherwise }
\end{array}\right.
$$

## State transition function

$$
\mathrm{M}_{+}=\left(\mathbb{N}, 0, \rho_{+}\right)
$$

$$
\rho_{+}(\sigma, j, \delta)= \begin{cases}\sigma+1 & \text { if } \delta=\downarrow \wedge \\ & \exists \mathcal{E}, e_{1}, e_{2}, v: j=\mathcal{E} \vdash e_{1}+e_{2} \Rightarrow v \\ \sigma & \text { otherwise }\end{cases}
$$

(State $X$ Judgement $X$ Direction) $\rightarrow$ State

## State transition function

$$
\mathrm{M}_{+}=\left(\mathbb{N}, 0, \rho_{+}\right)
$$

$$
\rho_{+}(\sigma, \mathcal{J}, \delta)= \begin{cases}\sigma+1 & \text { if } \delta=\downarrow \wedge \\ \sigma & \exists \mathcal{E}, e_{1}, e_{2}, v: j=\mathcal{E} \vdash e_{1}+e_{2} \Rightarrow v \\ \sigma & \text { otherwise }\end{cases}
$$

(State $X$ Judgement $X$ Direction) $\rightarrow$ State

## State transition function

$$
\mathrm{M}_{+}=\left(\mathbb{N}, 0, \rho_{+}\right)
$$

$$
\rho_{+}(\sigma, \mathcal{J}, \mathcal{\delta})= \begin{cases}\sigma+1 & \text { if } \delta=\downarrow \wedge \\ & \exists \mathcal{E}, e_{1}, e_{2}, v: j=\mathcal{E} \vdash e_{1}+e_{2} \Rightarrow v \\ \sigma & \text { otherwise }\end{cases}
$$

(State $X$ Judgement $X$ Direction) $\rightarrow$ State

## State transition function

$$
\mathrm{M}_{+}=\left(\mathbb{N}, 0, \rho_{+}\right)
$$

$$
\rho_{+}(\sigma, \mathcal{J}, \mathcal{\delta})= \begin{cases}\sigma+1 & \begin{array}{c}
\text { if } \delta=\downarrow \wedge \\
\exists \mathcal{E}, e_{1}, e_{2}, v: j=\mathcal{E} \vdash e_{1}+e_{2} \Rightarrow v \\
\text { otherwise }
\end{array}\end{cases}
$$

(State $X$ Judgement $X$ Direction) $\rightarrow$ State

## Monitoring with $\rho_{+}$

$$
\begin{array}{cc}
\hline \sigma_{2}^{\uparrow} \uparrow \emptyset \vdash 2 \Rightarrow 2 \downarrow \sigma_{2}^{\downarrow} & \\
\sigma_{3}^{\uparrow} \uparrow \emptyset \vdash 1 \Rightarrow 1 \downarrow \sigma_{3}^{\downarrow} \\
& \sigma_{1}^{\uparrow} \uparrow \emptyset \vdash 1+2 \Rightarrow 3 \downarrow \sigma_{1}^{\downarrow}
\end{array}
$$

$$
\rho_{+}(0,(\emptyset \vdash 1+2 \Rightarrow 3), \uparrow)
$$

## Monitoring with $\rho_{+}$

$$
\begin{array}{ccc}
\overline{\sigma_{2}^{\uparrow} \uparrow \emptyset \vdash 2 \Rightarrow 2 \downarrow \sigma_{2}^{\downarrow}} & \overline{\sigma_{3}^{\uparrow} \uparrow \emptyset \vdash 1 \Rightarrow 1 \downarrow \sigma_{3}^{\downarrow}} \quad & \sigma_{4}^{\uparrow \uparrow \emptyset \vdash+\Rightarrow+\downarrow \sigma_{4}^{\downarrow}} \\
\sigma_{1}^{\uparrow \uparrow \emptyset \vdash 1+2 \Rightarrow 3 \downarrow \sigma_{1}^{\downarrow}}
\end{array}
$$

$$
\rho_{+}(0,(\emptyset \vdash 1+2 \Rightarrow 3), \uparrow)
$$

## Monitoring with $\rho_{+}$

$$
\rho_{+}(0,(\emptyset \vdash 1+2 \Rightarrow 3), \uparrow)
$$

## Monitoring with $\rho_{+}$

$$
\begin{aligned}
\hline \sigma_{2}^{\uparrow} \uparrow \emptyset \vdash 2 \Rightarrow 2 \downarrow \sigma_{2}^{\downarrow} & \sigma_{3}^{\uparrow} \uparrow \emptyset \vdash 1 \Rightarrow 1 \downarrow \sigma_{3}^{\downarrow}
\end{aligned} \cdots \sigma_{4}^{\uparrow \uparrow \emptyset \vdash+\Rightarrow+\downarrow \sigma_{4}^{\downarrow}}
$$

## Monitoring with $\rho_{+}$

$\overline{(0) \uparrow \emptyset \vdash 2 \Rightarrow 2 \downarrow \sigma_{2}^{\downarrow}} \quad \overline{\sigma_{3}^{\uparrow} \uparrow \emptyset \vdash 1 \Rightarrow 1 \downarrow \sigma_{3}^{\downarrow}} \quad \overline{\sigma_{4}^{\uparrow} \uparrow \emptyset \vdash+\Rightarrow+\downarrow \sigma_{4}^{\downarrow}}$


## Monitoring with $\rho_{+}$



## Monitoring with $\rho_{+}$

$\frac{\overline{0 \uparrow \emptyset \vdash 2 \Rightarrow 2 \downarrow 0} \quad \frac{\overline{0 \uparrow \emptyset \vdash 1 \Rightarrow 1 \downarrow 0}}{0 \uparrow \emptyset \vdash 1+2 \Rightarrow 3 \downarrow 1} \overline{0 \uparrow \emptyset \vdash+\Rightarrow+\downarrow 0}}{\rho_{+}(0,(\emptyset \vdash 1+2 \Rightarrow 3), \downarrow)}$

## Our work

- Monitoring for evaluation trees


## Product construction

## Evaluation

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- Monitoring for evaluation trees
- Product construction


## Evaluation

## Product construction

## User program

$$
\begin{aligned}
& \text { let } \mathrm{f} x= \\
& \text { if } x<0 \text { then } 0 \text { else } f(x-I)
\end{aligned}
$$

## Product construction

## User program

$$
\begin{aligned}
& \text { let } \mathrm{f} x= \\
& \text { if } x<0 \text { then } 0 \text { else } f(x-1)
\end{aligned}
$$

Monitoring code

## Product construction

## User program

Monitoring code

$$
\begin{aligned}
& \text { let } \mathrm{f} x= \\
& \text { if } \mathrm{x}<0 \text { then } 0 \text { else } \mathrm{f}(\mathrm{x}-\mathrm{I})
\end{aligned}
$$

$$
\operatorname{assert}(x>0)
$$

$$
\text { let } \mathrm{f} x=\text { if } \mathrm{x}<0 \text { then } 0 \text { else assert }(\mathrm{x}>0) \text {; } \mathrm{f}(\mathrm{x}-\mathrm{I})
$$

# Monitor specification 

$$
\mathrm{S}_{+}=(0, \text { mtrans_sum })
$$

# Monitor specification 

$$
\mathrm{S}_{+}=(0, \text { mtrans_sum })
$$

Initial state expression

## Monitor specification

$$
\mathrm{S}_{+}=(0, \text { mtrans_sum })
$$

Initial state expression

## State transformer procedure

## State transformer

## procedure

$$
\mathrm{S}_{+}=(0, \text { mtrans_sum })
$$

function mtrans_plus counter expression direction $=$ match expression, direction with
| <expr_patt< + \$_\$ >>, Down -> <expr_term< \$counter\$ + 1 >>
| _ -> counter

## State transformer

## procedure

$$
S_{+}=(0, \text { mtrans_sum })
$$

function mtrans_plus counter expression direction $=$ match expression, direction with
| <expr_patt< + \$_\$ >>, Down -> <expr_term< \$counter\$ + 1 >>
| _ -> counter

## (State X Expression X Direction) $\longrightarrow$ Monitoring code

## State transformer

## procedure

$$
S_{+}=(0, \text { mtrans_sum })
$$

function mtrans_plus counter expression direction $=$ match expression, direction with
| <expr_patt< + \$_\$ >>, Down ->
<expr_term< \$counter\$ + 1 >>
| _ -> Counter
(State $X$ Expression $X$ Direction) $\longrightarrow$ Monitoring code

## Product of $\mathrm{I}+2$ and $\mathrm{M}_{+}$

## $1+2$

## Product of $I+2$ and $M_{+}$

let $x$ _plus = (fun x_11 x_12 in
let x_1_plus_2
x_1_plus_2

$$
->x_{-} 11+x_{-} 12
$$

$$
=x_{-} p l u s 12 \text { in }
$$

## Product of $\mathrm{I}+2$ and $\mathrm{M}_{+}$

let $x$ _plus =
(fun x_11 x_12 s_1_pre -> x_11 + x_12, s_1_pre) in
let x_1_plus_2, s_1_plus_2 = x_plus 12 s in x_1_plus_2, s_1_plus_2 + 1

## High order program

let rec fold_left f accu l = match l with
| [] -> accu
| a::t ->

$$
\text { let accu' } \quad=f \text { accu } a
$$

in
fold_left f accu' t

## High order product

let rec fold_left_m f accu l c = match l with

$$
\begin{aligned}
& \text { | [] -> accu, c } \\
& \text { a::t -> }
\end{aligned}
$$

let accu', c' = f accu a (c + 1) in fold_left_m f accu' t c'

## High order product

let rec fold_left_m f ecu lc = match l with

$$
\begin{aligned}
& \text { | [] -> accu, c } \\
& \text { | a::t -> }
\end{aligned}
$$

let ecu', $c^{\prime}=f$ ecu a ( $\left.c+1\right)$ in fold_left_m f ecu' tc'

## High order product

let rec fold_left_m f accu l c = match l with

$$
\begin{aligned}
& \text { | [] -> accu, c } \\
& \text { a: :t -> } \\
& \text { let accu', } c^{\prime}=f \text { accu a }(c+1) \text { in } \\
& \text { fold_left_m } f \text { accu, } t \text { c, }
\end{aligned}
$$

## Our work

- Monitoring for evaluation trees
- Product construction


## Evaluation

## Our work

- Monitoring for evaluation trees
- Product construction
- Evaluation


## FunV

# - Product construction algorithm 

- Reachability checker Dsolve [1]
[1] M. Kawaguchi, P. M. Rondon, and R. Jhala. Type-based data structure verification. In PLDI, 2009.


## Evaluation

## - 600+ LOC

- 62 benchmarks [1-4]


Figure 8: Experiments

## Evaluation

## - 600+ LOC

- 62 benchmarks [1-4]
[1] M. Hofmann. A type system for bounded space and functional in-place update-extended abstract. In $E S O P, 2000$.
[2] M. Hofmann. The strength of non-size increasing computation. In POPL, 2002.
[3] M. Hofmann and S. Jost. Static prediction of heap space usage for firstorder functional programs. In POPL, 2003.
[4] N. Kobayashi. Types and higher-order recursion schemes for verification of higher-order programs. In POPL, 2009.


Figure 8: Experiments.

## Evaluation: benchmarks

- Recursion
- Higher-order functions
- Algebraic data types
- Abstract data types



## Evaluation: properties

- Safety and termination
- Linear inequalities over values and measures
- Inclusion checks over sets of program values
- Ranking function/transition invariant checks



## Demo

## Conclusion

- Monitoring
- Product construction
- Evaluation

